

# CISC 839

## Software Engineering of Usable Computing Systems

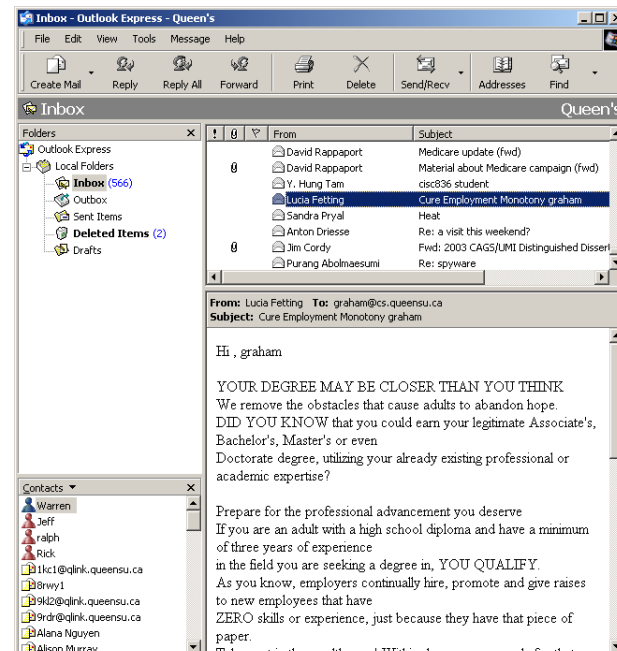
**T.C. Nicholas Graham**

<http://www.cs.queensu.ca/~graham/cisc836>

## Implementing Groupware

## Examples of Asynchronous Groupware Products

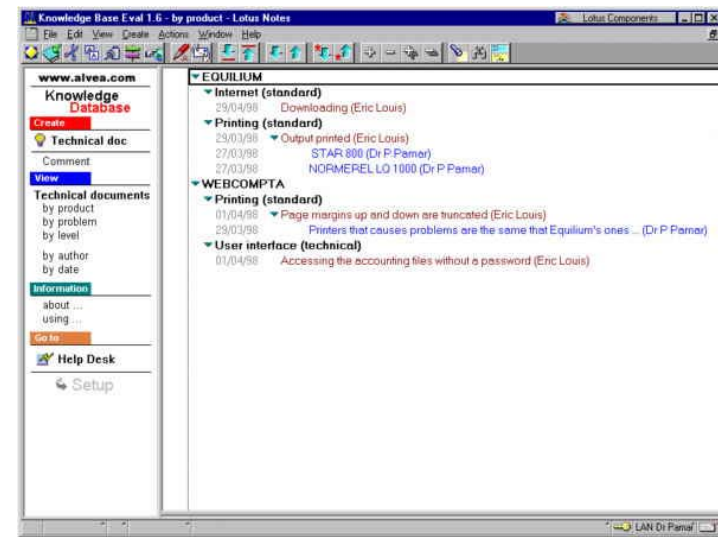
- ⌘ Email
- ⌘ Discussion Forum
- ⌘ Shared document space
- ⌘ Shared repository



Outlook  
Express



## WebCT Forum



## Lotus Notes

[http://www.alvea.com/ehhtml/know\\_screen\\_shot.htm](http://www.alvea.com/ehhtml/know_screen_shot.htm)

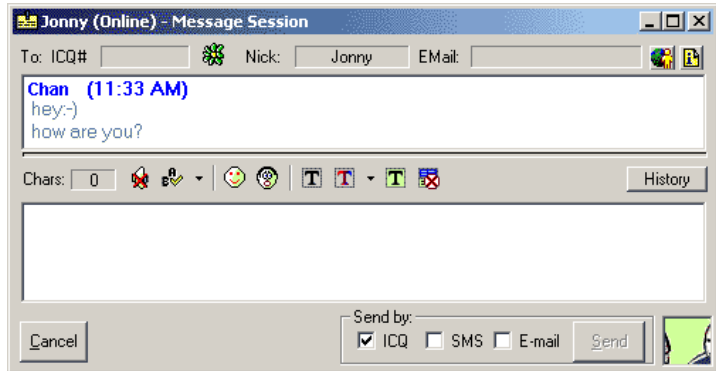


## Bugzilla

## Examples of Synchronous Groupware

- ⌘ Instant messaging
- ⌘ Multiuser video games
- ⌘ Electronic meeting systems

# ICQ Instant Messaging

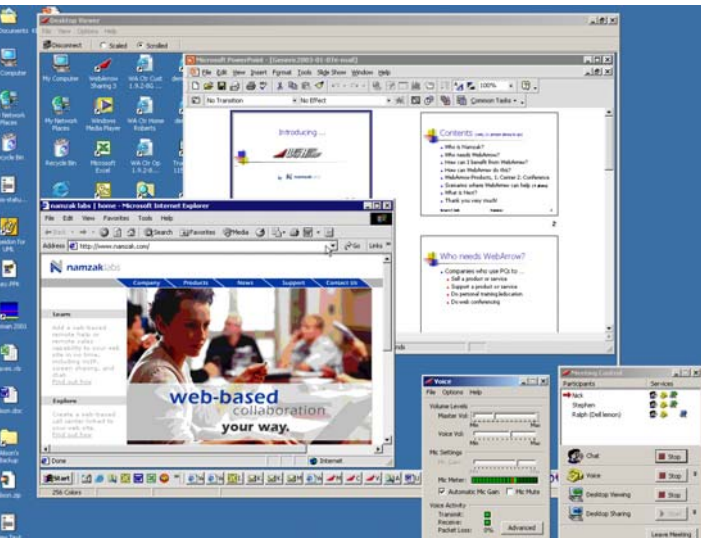


<http://www.icq.com/icqtour/simple/user-menu.html>

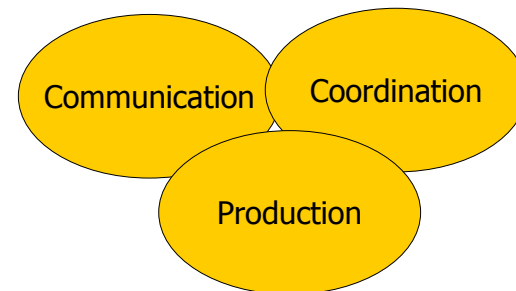
# Dark Age of Camelot Multiuser Video Game



# WebArrow/Conference Electronic Meeting System



# Clover Model of Communication [Calvary et al., 1997]

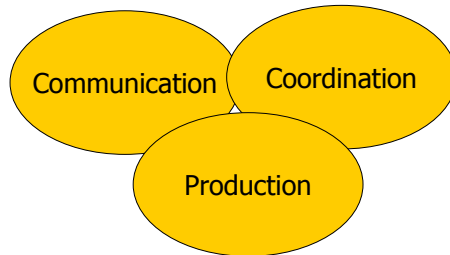


# Clover Model

⌘ People convey information to each other in three ways (termed *spaces*)

☒ *Communication*: direct exchange of information, e.g., speech, gestures, facial expression, gaze information

☒ Supported by (e.g.) telephone, videophone

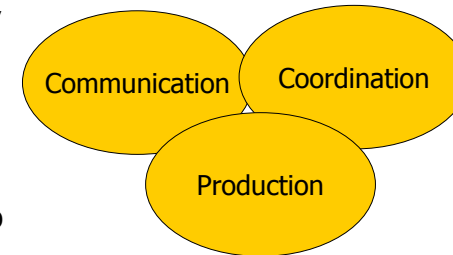


# Clover Model

⌘ *Production*: exchange of information through some artifact being collaboratively produced

⌘ E.g., shared whiteboard, shared text editor

⌘ Supported by (e.g.) MS NetMeeting, multiuser Video Annotator, Ultima Online



# Clover Model

⌘ *Coordination*: organization of how collaboration is to be enacted. E.g.,

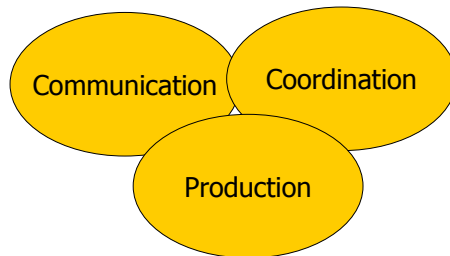
☒ In a chess game, use turn-taking

☒ For patient's medical records, use workflow

☒ For software development, use locking at class level

☒ For collaborative drawing, use social protocols

⌘ Supported by workflow systems, concurrency control algorithms

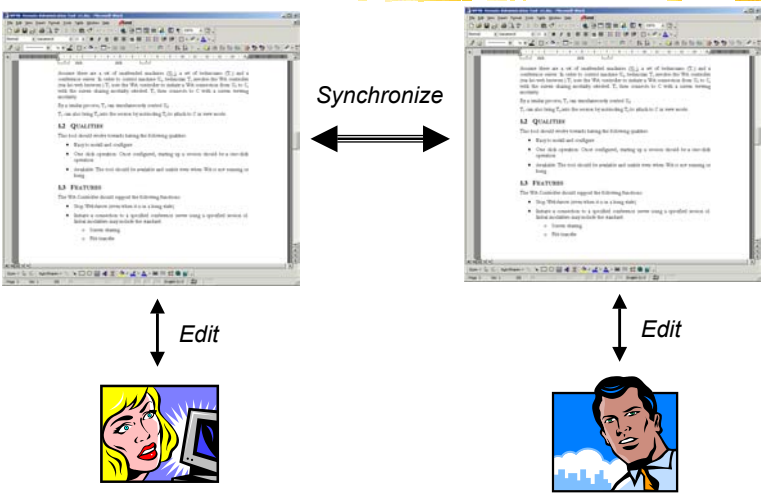


# Concurrency Control

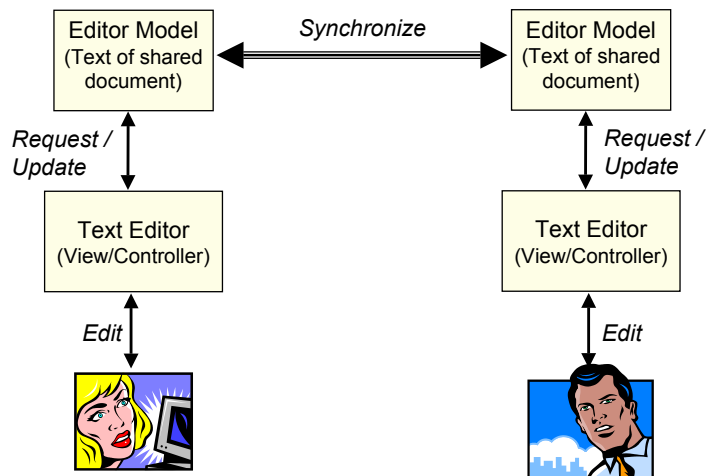
# Problem

- ⌘ Two people simultaneously update the same shared context
- ⌘ Updates may conflict
- ⌘ Concurrency control: algorithms that ensure:
  - ☑ All users' views of the shared context agree
  - ☑ In case of conflicts, result is intuitive

# Example: Shared Text Editor



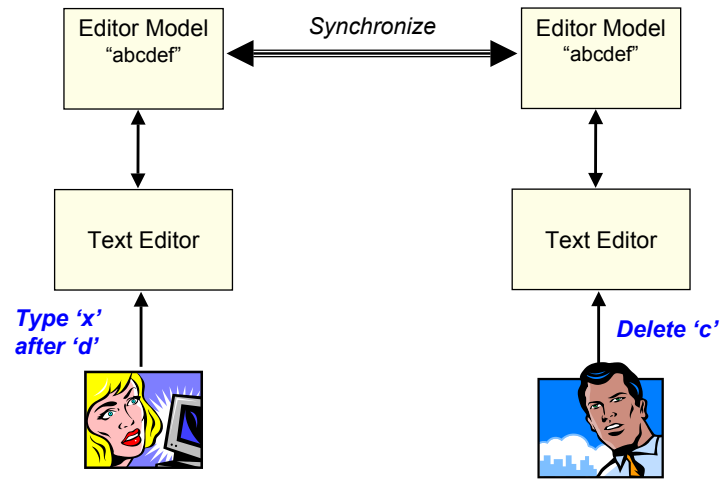
# Example: Shared Text Editor



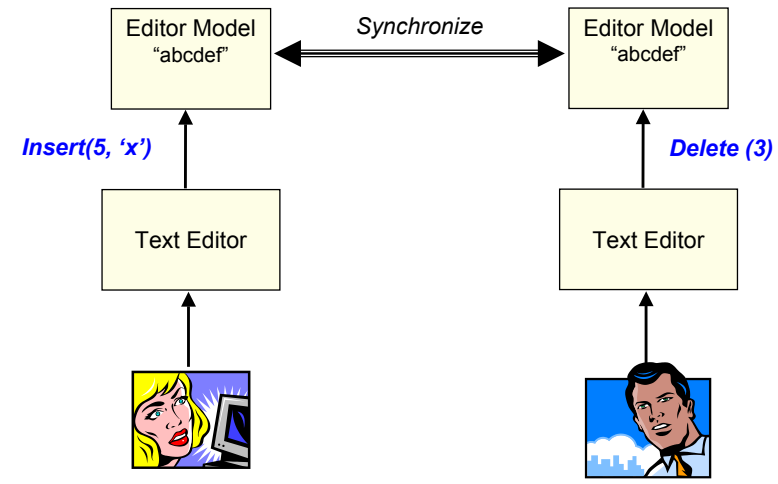
# Editor Model

- ⌘ Insert (int pos, char c)
  - ☑ Inserts  $c$  before position  $pos$ , where  $pos \geq 1$
- ⌘ Delete (int pos)
  - ☑ Deletes character at position  $pos$ , where  $pos \geq 1$

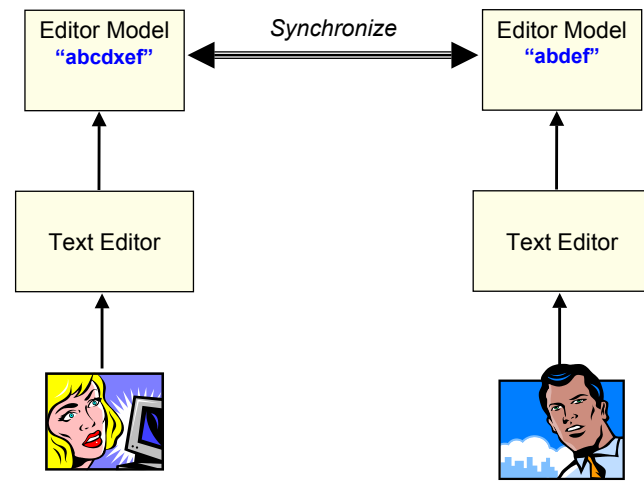
# Example: Shared Text Editor



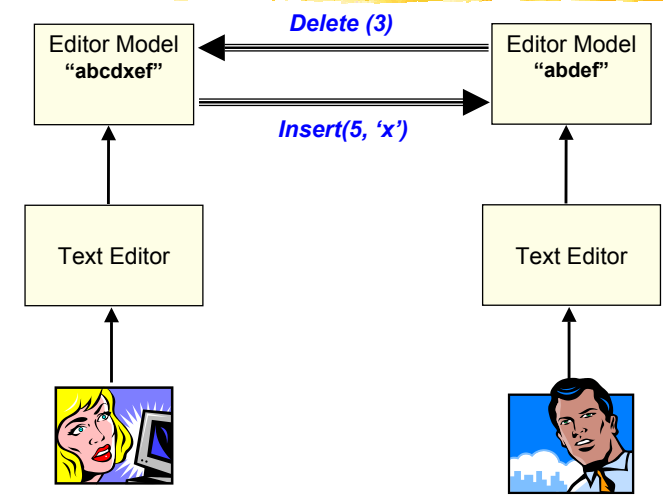
# Example: Shared Text Editor



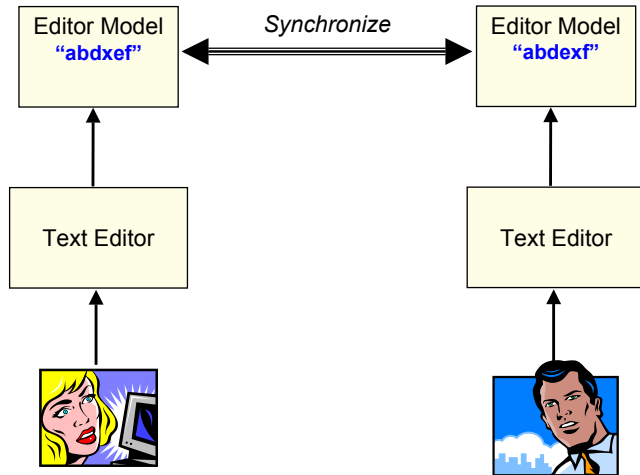
# Example: Shared Text Editor



# Example: Shared Text Editor



## Example: Shared Text Editor



## Problem

- ⌘ Replicated data
- ⌘ Performing operations in different order results in divergent values in text editor model
  - ☒ Users will see different documents on their screen

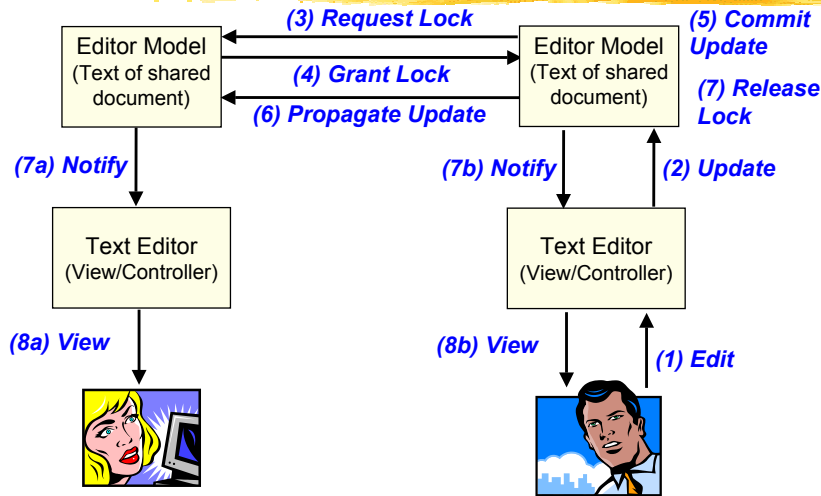
## Goals of Concurrency Control

- ⌘ Maintain consistent versions of models for all users
- ⌘ Provide intuitive result when conflicts occur
  - ☒ E.g., could solve conflicts in text editor by setting document to "". Not an intuitive result!
- ⌘ Maintain adequate feedthrough, feedback times
  - ☒ Many algorithms make feedback times unacceptable

## Locking

- ⌘ Before performing update to shared context, must obtain a lock
- ⌘ Negotiated amongst all instances of shared context
  - ☒ If another user holds the lock, must wait until it is released
- ⌘ Release lock when update has been propagated through system

## Example: Locking



## Locking

- ⌘ Simple to implement
- ⌘ Intuitive results
  - ☑ No conflicts
- ⌘ Poor feedback time
  - ☑ Must go over network to obtain lock before user sees result of change
  - ☑ Bad for interactive tasks like text editing

## Rollback

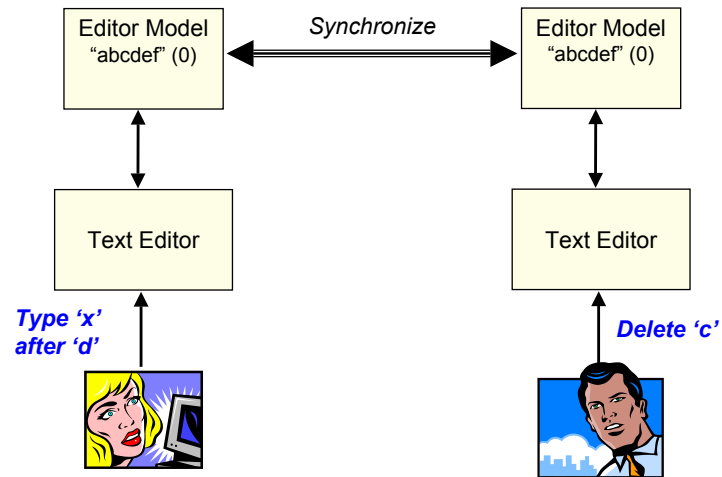
- ⌘ Detect conflicts via timestamp
- ⌘ When conflicts occur
  - ☑ Roll back to earlier state
  - ☑ Reapply updates in canonical order
    - ☑ I.e., same order on each machine

## Timestamps

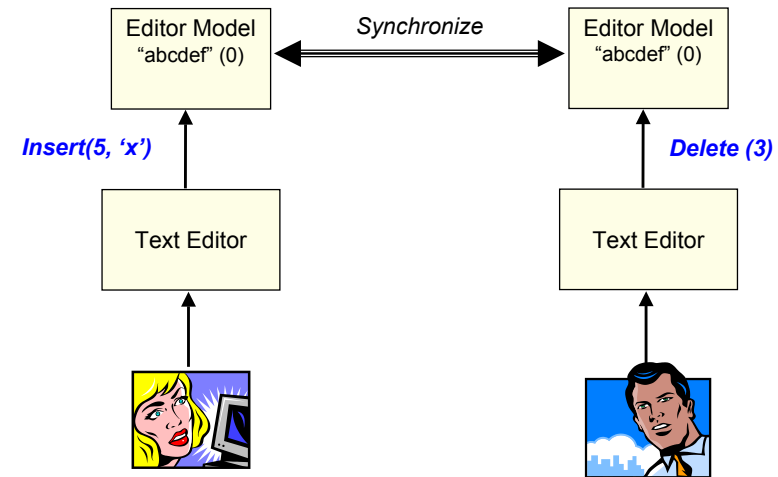
- ⌘ Simply increment a counter every time an update is performed
  - ☑ Counter is timestamp
- ⌘ Include timestamp in messages carrying remote updates
  - ☑ Specifies version of model against which update applied
- ⌘ If remote update has older timestamp than model, a conflict has occurred



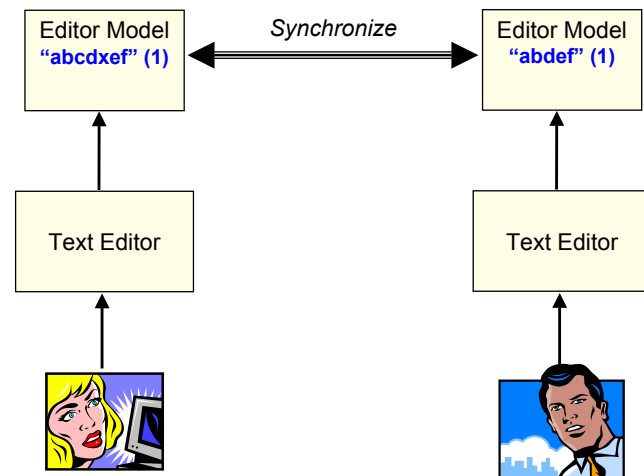
# Example: Rollback



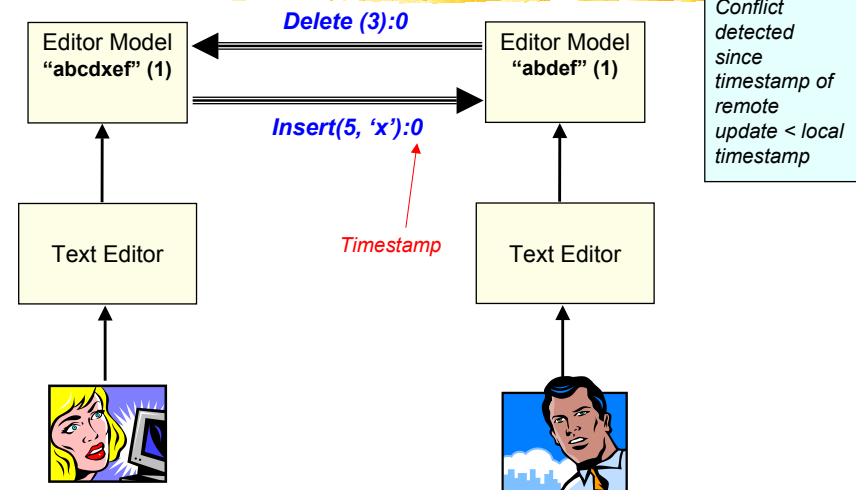
# Example: Rollback



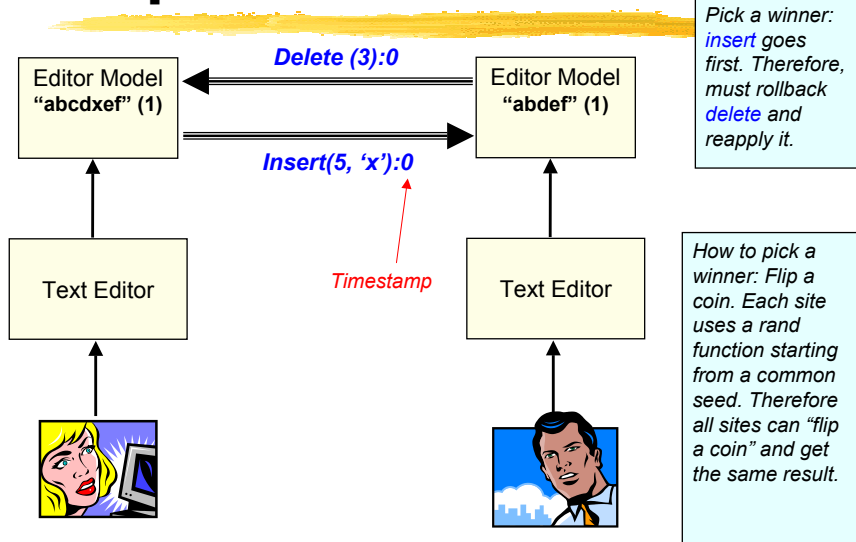
# Example: Rollback



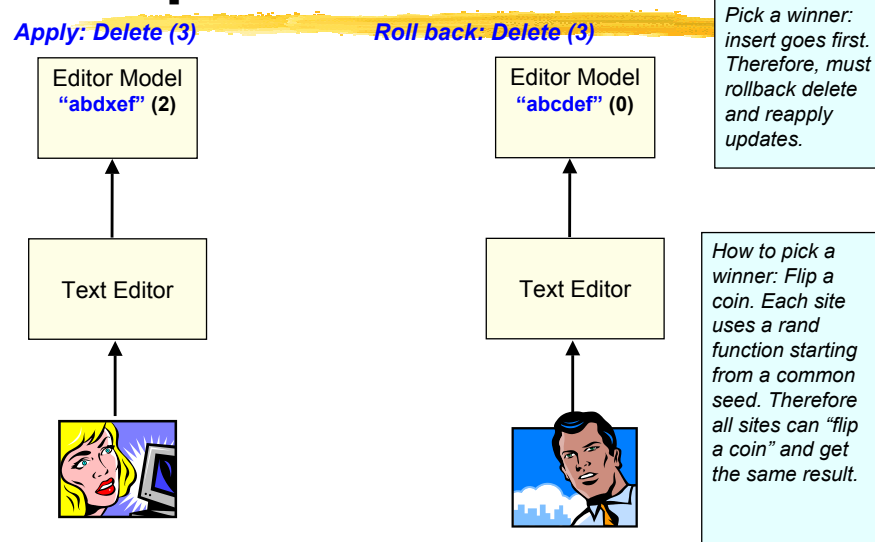
# Example: Rollback



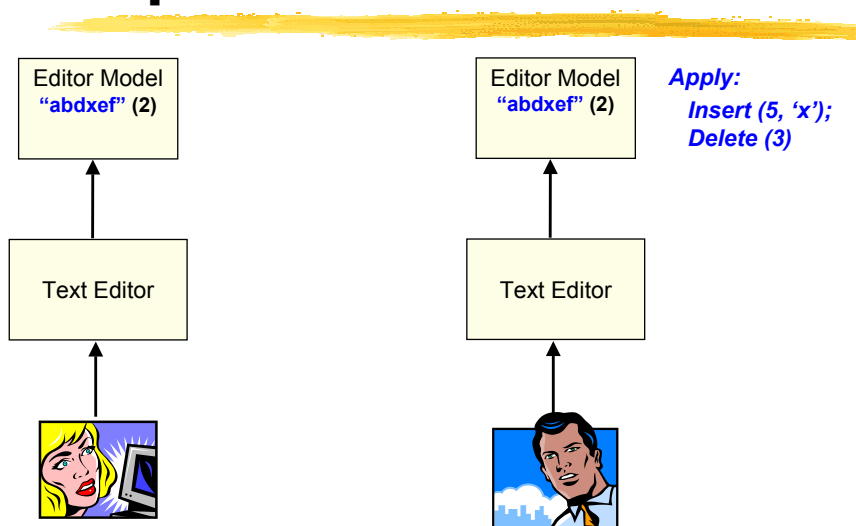
# Example: Rollback



# Example: Rollback



# Example: Rollback



# Rollback

- ⌘ Immediate feedback
- ⌘ Requires ability to roll back all operations
  - ☑ May be hard to program, costly in resources
  - ☑ Some operations can't be rolled back – e.g., file save, send email, "launch" button on nuclear missile
- ⌘ Rollbacks may be confusing to user
  - ☑ See result of input, then it's changed later

# Operation Transform

⌘ Similar to rollback, except:

- ☒ When conflict detected, instead of rolling back, *transform* operations
- ☒ Requires *fixup* operations that undo and reapply all in one operation

# Operation Transforms

⌘ Example: Should have executed:

☒ Insert (5, 'x') ; Delete (3)     *"abcdef" → "abdxef"*

⌘ Actually executed:

☒ Delete (3)     *"abdef"*

⌘ Now want to execute:

☒ Delete (3)<sup>-1</sup> ; Insert (5, 'x') ; Delete (3)

☒ ≡ Insert (4, 'x') ; Delete (3)<sup>-1</sup> ; Delete (3)

☒ ≡ Insert (4, 'x')     *"abdxef"*

# Operation Transforms

⌘ No need to roll back

☒ ⇒ No need to save earlier states

⌘ Need to provide "fix up" operations for pairs of operations in model

- ☒  $n^2$  operations
- ☒ Some may be hard to implement
- ☒ Some may not have intuitive fix ups

# Concurrency Control

⌘ Three methods:

- ☒ Locking
- ☒ Rollbacks
- ☒ Operation Transform

⌘ See next week's readings for more types

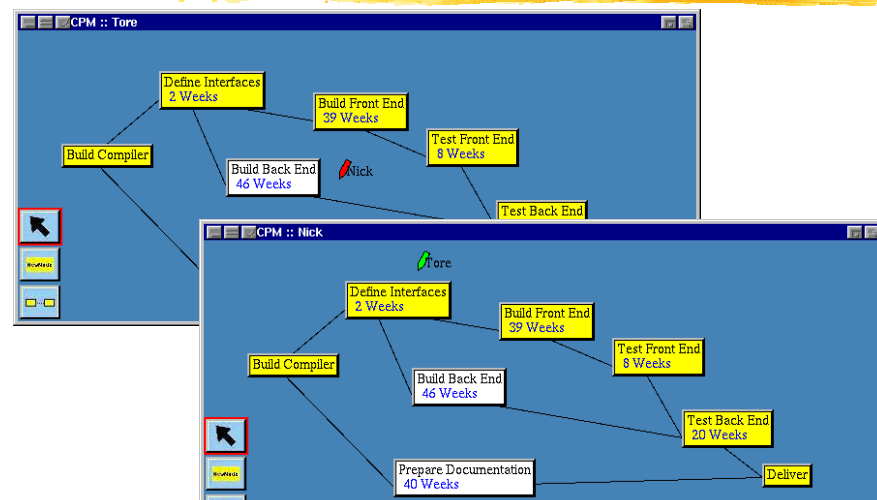
# Groupware Architectures

- ⌘ Clock
- ⌘ Dragonfly
- ⌘ ...plus see next week's readings

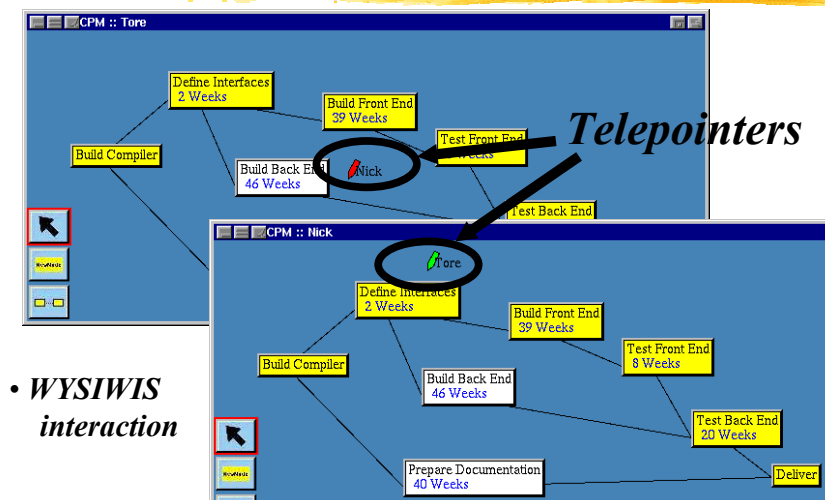
# Outline

- ⌘ Why is developing groupware hard?
- ⌘ Programming abstractions for groupware
  - ⌘ Declarative development of groupware
  - ⌘ Conceptual architectures as programs
- ⌘ Flexible implementation of groupware
  - ⌘ Annotations for implementation
  - ⌘ The Dragonfly implementation architecture
- ⌘ Timings

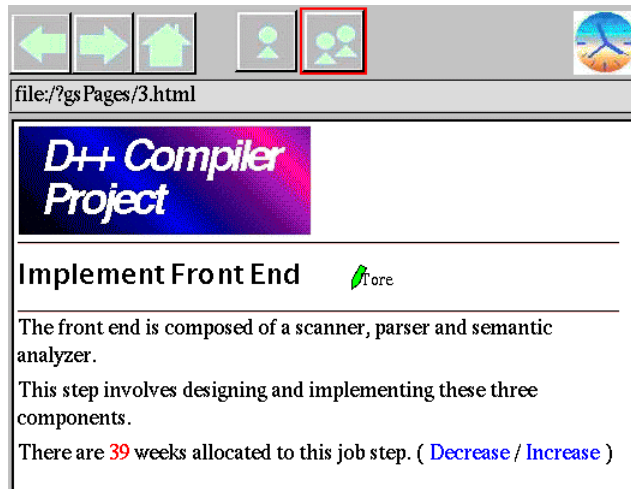
## Synchronous Critical Path Scheduling Application



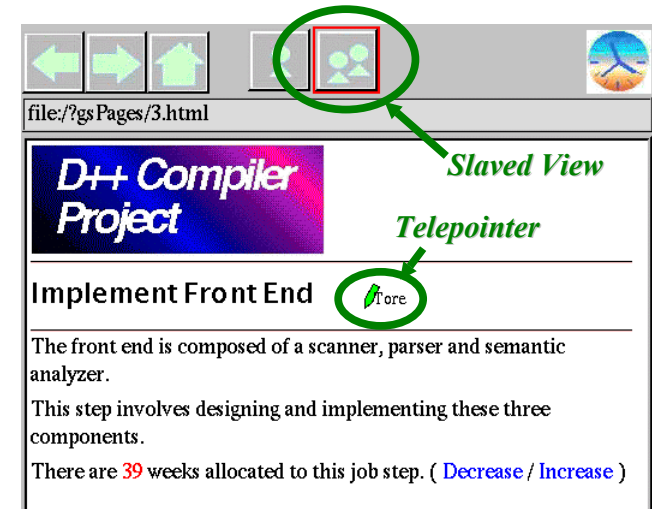
## Synchronous Critical Path Scheduling Application



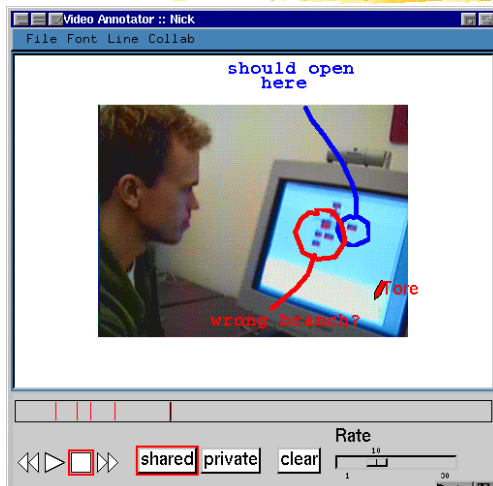
# GroupScape Web Browser



# GroupScape Web Browser



# Synchronous Video Annotator



Two users

- Annotating tape of usability testing session
- Standard VCR controls
- Annotations saved with video frame

# Why is Developing GW Hard?

- ⌘ Synchronizing Relaxed WYSIWIS Views
- ⌘ Concurrency
  - ⌘ Race Conditions
- ⌘ Distribution
  - ⌘ Networking, Latency, Scalability, ...
- ⌘ Multimedia
  - ⌘ Synchronizing, Integrating MM Streams
- ⌘ User-Centered Design

# Outline

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# Conceptual Architectures

*High level architecture style hides:*

- ⌘ Distribution
- ⌘ Networking
- ⌘ Concurrency control

*...allowing developer to concentrate on  
functionality of application*

- ⌘ E.g., Clock, PAC, ALV, C2, ...

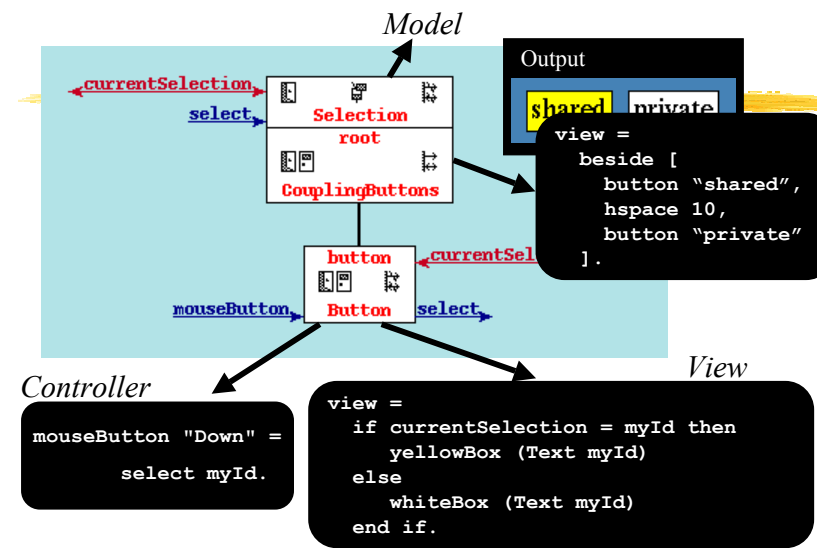
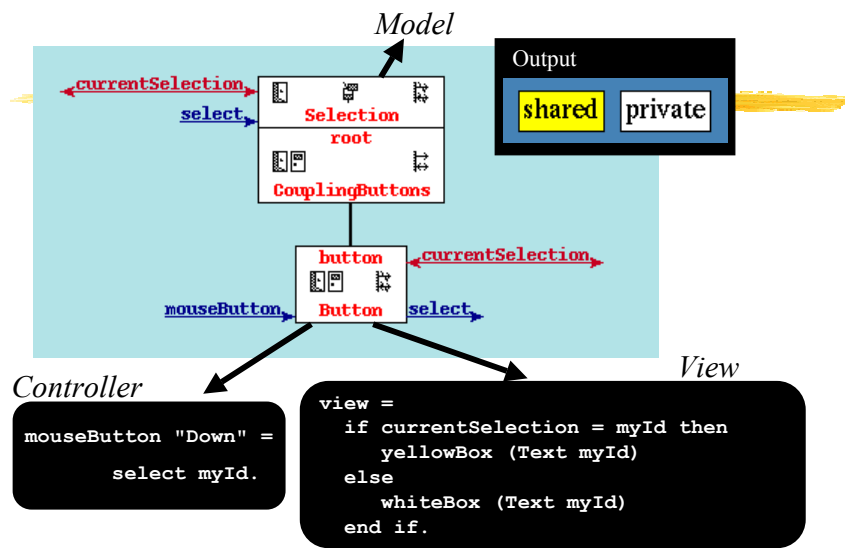
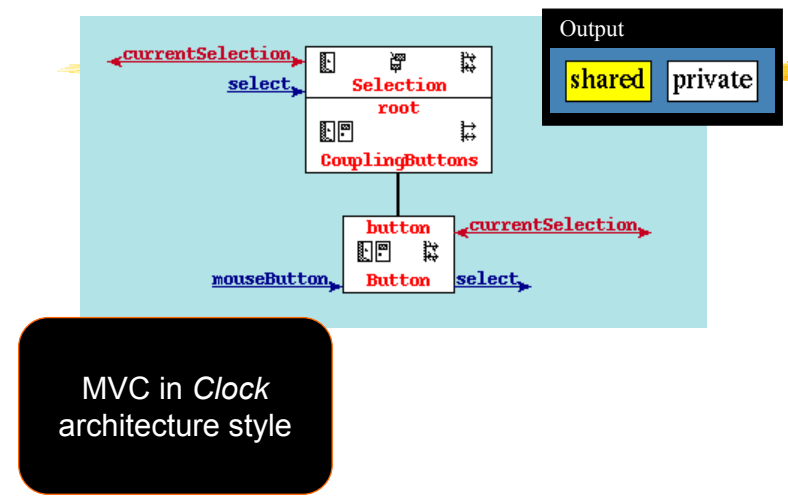
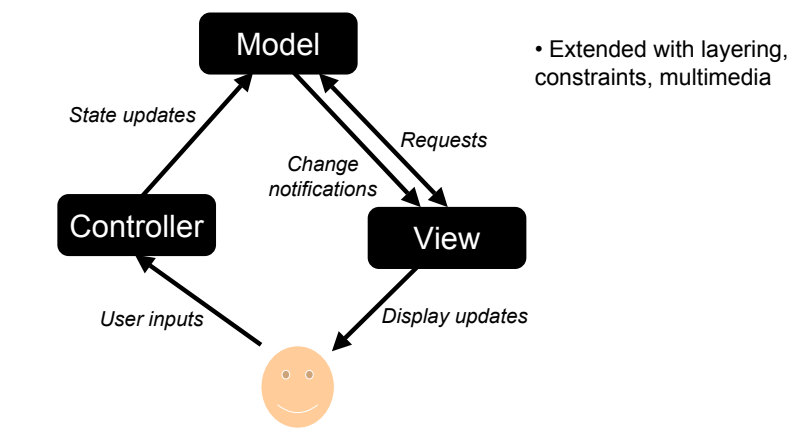
# Clock Architecture Style

- ⌘ Conceptual architecture that is also a program
- ⌘ ClockWorks visual programming environment
- ⌘ Connectors encode temporal properties
  - ☒ Concurrency control
  - ☒ Temporal media

# Clock Architecture Style

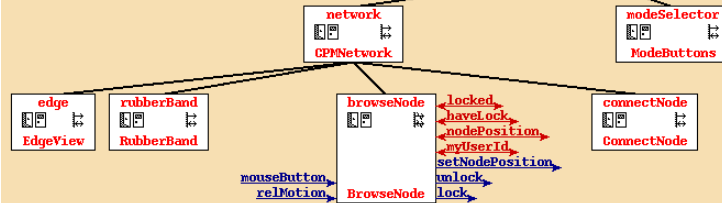
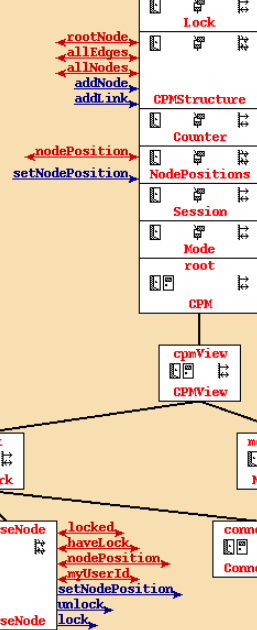
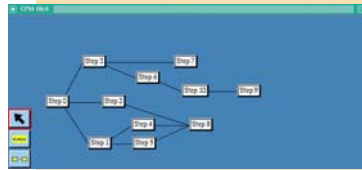
- ⌘ Two views
  - ☒ Layered MVC
  - ☒ PAC where Control encoded by connectors

# MVC Architecture



## Layered Architectures in Clock

- Updates send information up tree
- Constraints pull information down

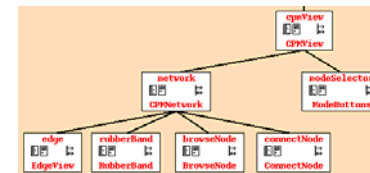


## Applying the Clock architecture to Groupware

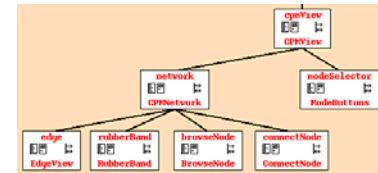
- Shared context at root of tree
- One instance of UI per user

setNodePosition "2"  
(200, 321)

Trigger view update  
for node 2

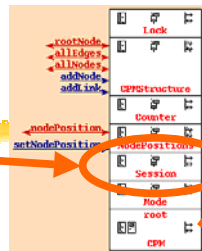


User 1



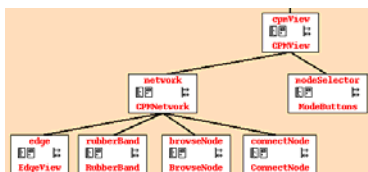
User 2

Information  
provided by  
session  
manager

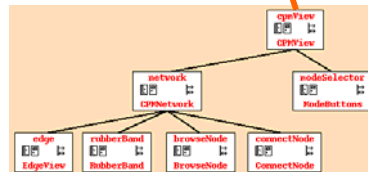


```
view =
  beside (
    map cpmView allUsers).
```

```
view =
  Window myUserId myHost (...).
```



User 1



User 2

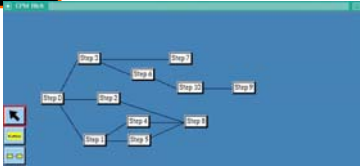
## Tree Structure

- ⌘ View composition
- ⌘ Data visibility
- ⌘ Concurrency control
  - ☑ Avoiding race conditions between user inputs
- ⌘ Synchronization of multimedia streams
  - ☑ Sharing sound, video, animations



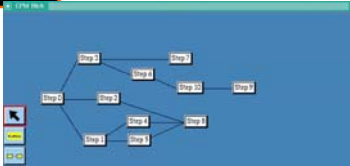
# Concurrency Control

```
% Create a new node
mouseButton "Down" =
  all [addNode getCount mousePosition,
       incrementCount].
```



# Concurrency Control

```
% Create a new node
mouseButton "Down" =
  all [addNode getCount mousePosition,
       incrementCount].
```

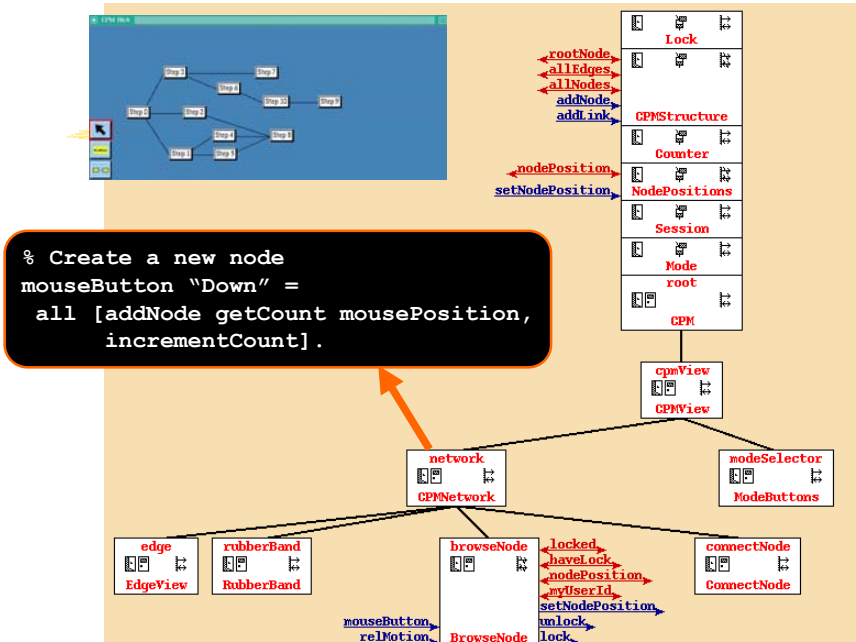


Possible execution order

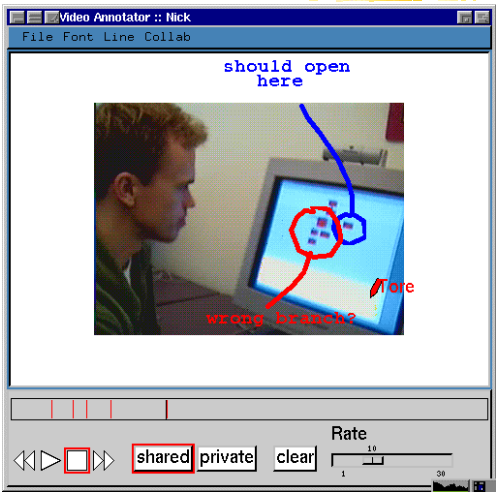
| User 1              | User 2              |
|---------------------|---------------------|
| getCount            |                     |
| addNode (... ,... ) | getCount            |
| incrementCount      | addNode (... ,... ) |
|                     | incrementCount      |

# Concurrency control

- ⌘ Each user's action leads to a *transaction*
  - 📦 request\* update
- ⌘ Transactions must be executed in a serializable fashion



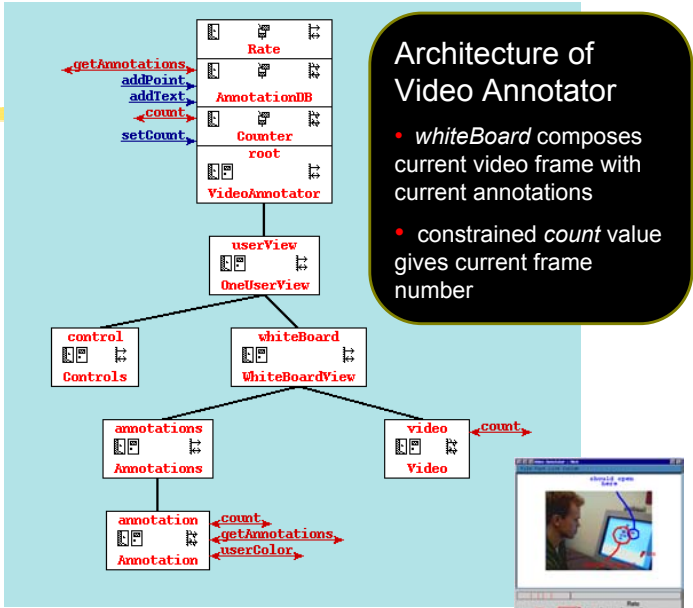
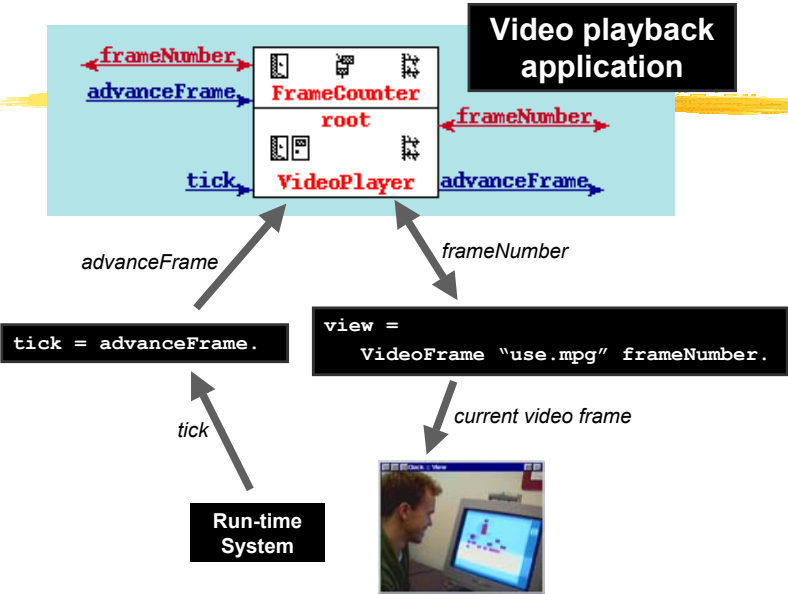
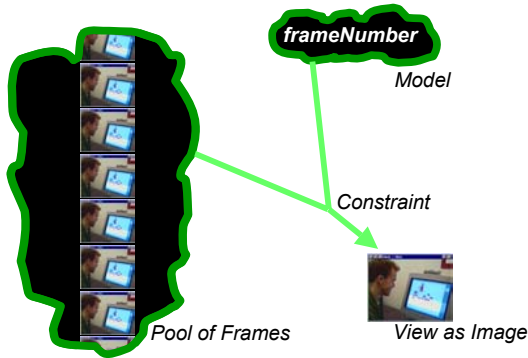
# Multimedia Groupware



- Synchronization of multimedia streams
- Integration of temporal, static media

## Extending MVC for Temporal Media

- *Constraints* remove time from temporal media
- *Composition* allows integration of temporal and static media



- Architecture of Video Annotator**
- *whiteBoard* composes current video frame with current annotations
  - constrained *count* value gives current frame number

## Clock Architecture Style: Summary

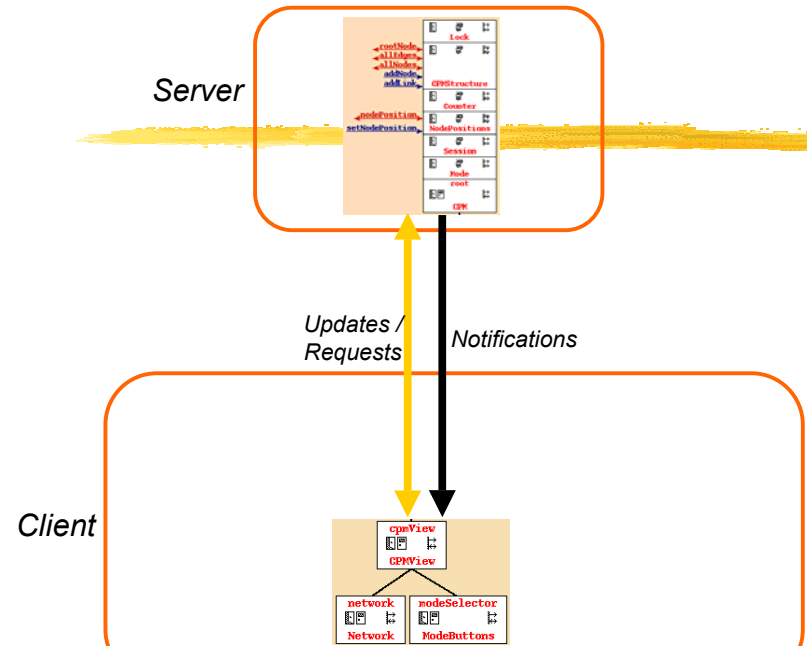
- ⌘ Conceptual architecture is part of program
- ⌘ Constraints help in:
  - ☑ Synchronizing relaxed WYSIWIS views
  - ☑ Multimedia
- ⌘ Concurrency control guaranteed by connectors
- ⌘ Visual editor for architectures

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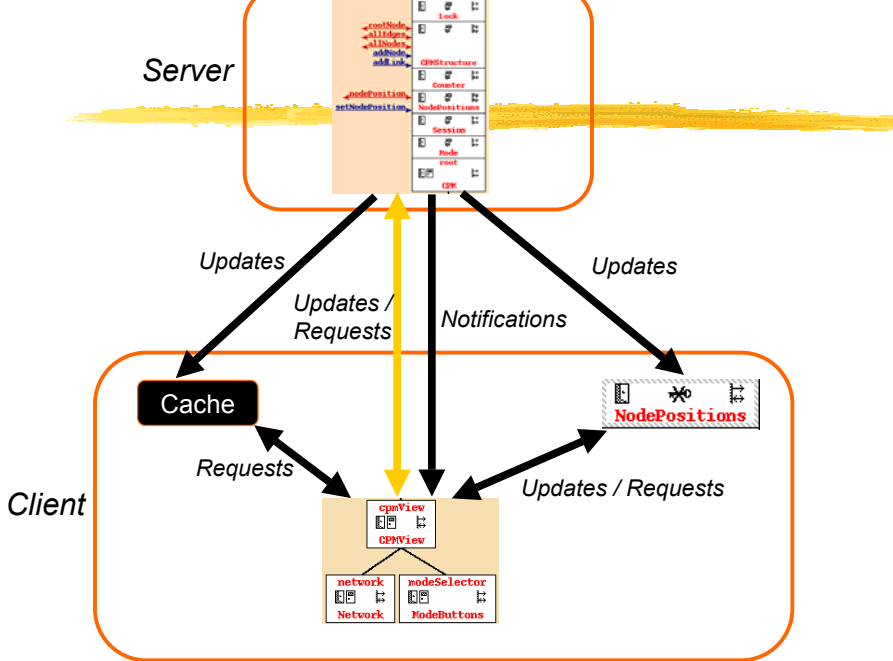
# Strategies for Cheating Latency

- ⌘ Sacrifice group response time
- ⌘ Reduce granularity of group actions
- ⌘ Take turns
- ⌘ Risk roll-backs
- ⌘ Buffer
- ⌘ Spend CPU, memory, bandwidth



# Space of Implementation Techniques

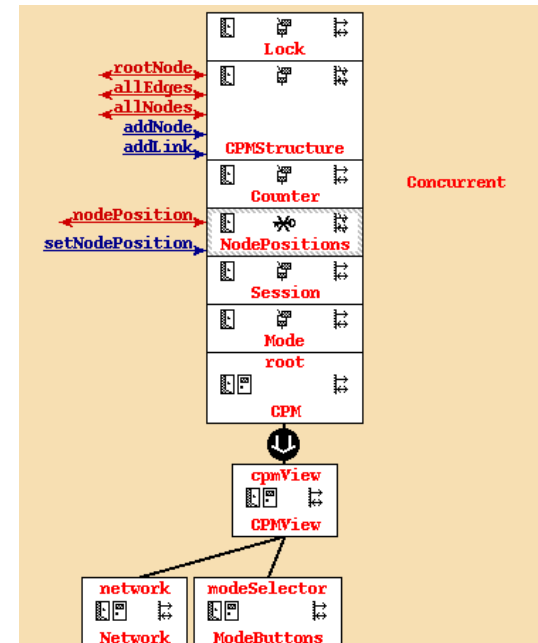
- ⌘ Concurrency control: *none, optimistic, pessimistic*
- ⌘ Topology: *tree, star, mesh, point*
- ⌘ Replication: *centralized, partially replicated, replicated*
- ⌘ Cache activity: *passive, active*
- ⌘ Cache location: *standard, mirror*
- ⌘ Network delivery: *lossless, lossy*
- ⌘ Network order: *FIFO, any*



## What Strategy is Best?

Depends on:

- ⌘ Application characteristics
  - ☐ Symmetry, roles, synchronicity
- ⌘ Hardware characteristics
  - ☐ How cheap is bandwidth, latency, CPU, memory?



### Annotations for Distributed Implementation

- Client / Server Split:



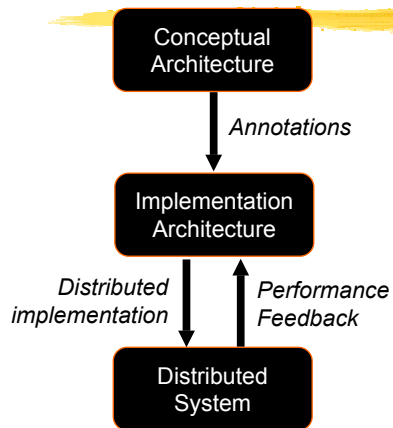
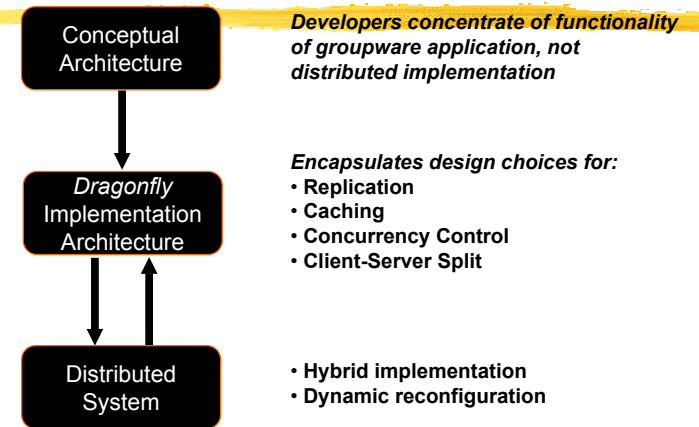
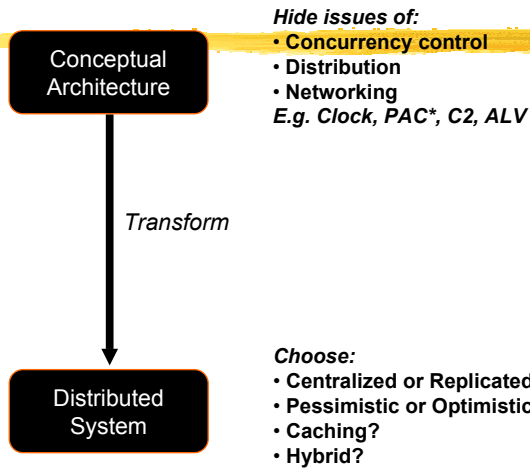
- Replication



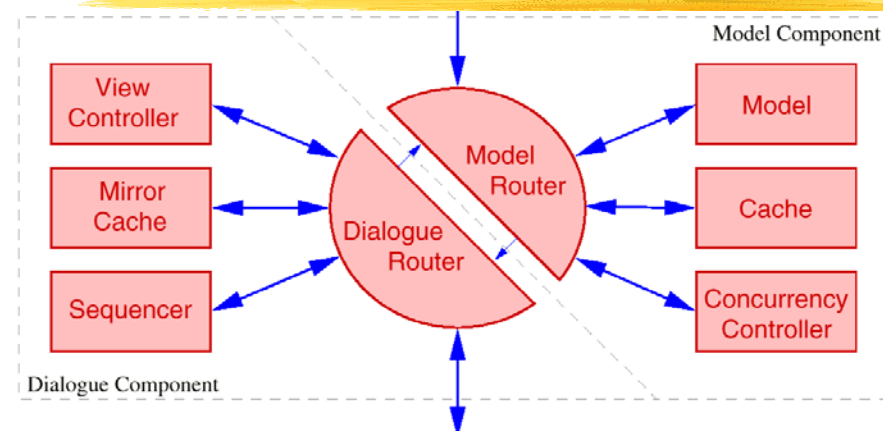
- Concurrency Control

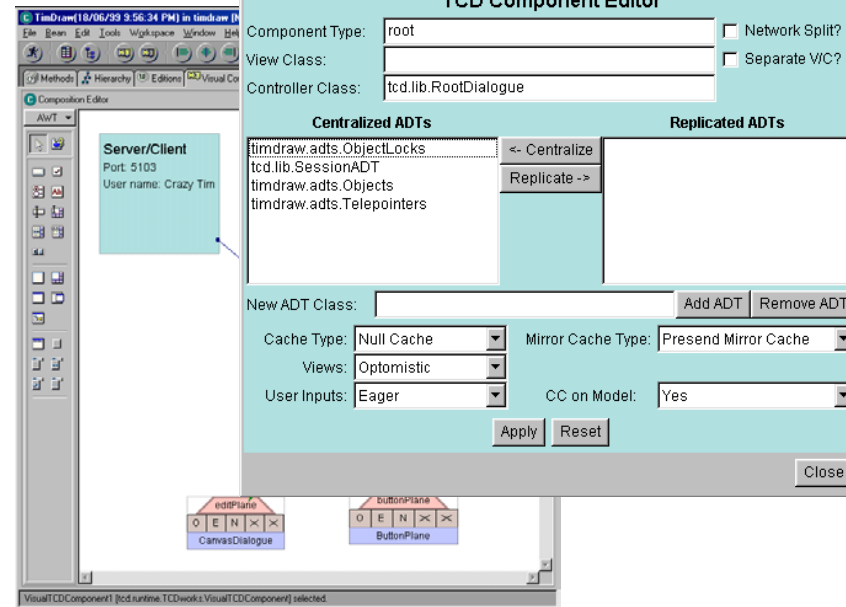
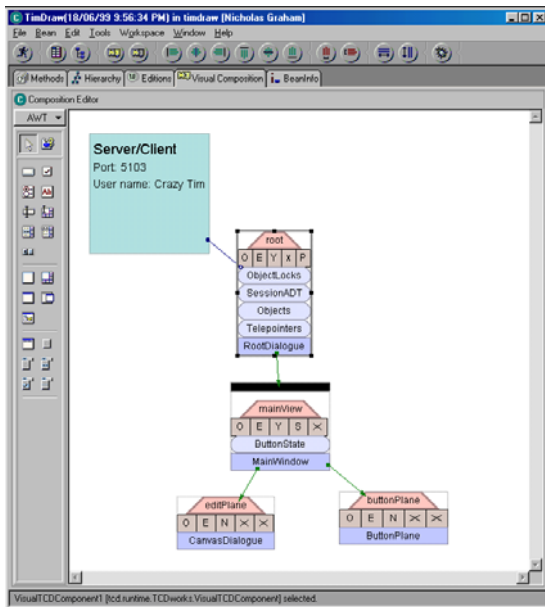
**Concurrent**

**Serialized**



## The Dragonfly Component

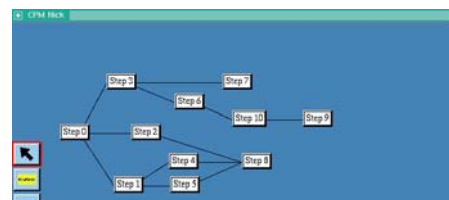




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# Timing Clock and Dragonfly



Response time where:

- Server: LAN: UltraSparc 1, York U.; WAN: 100 MHz R4000 SGI Indy, Queen's University
- Client: UltraSparc 1, York U.
- Network Latency: LAN: 1 ms; WAN: 11 ms

# Conclusions



- ⌘ Large gap between design-level architecture and distributed system
- ⌘ Dragonfly acts as layer between, allowing experimentation with implementation strategies
- ⌘ Programmers can give high-level guidance for implementation